

Static clusters in cellular Networks. The Nearest Neighbor Model.

Luis David Álvarez Corrales
with **Anastasios Giovanidis & Laurent Decreusefond.**

Télécom ParisTech, France.

4th Stochastic Geometry Days, 27.08.2015.

Introduction

- ▶ In recent years, it became popular to model the irregularity of the network base station (BS) locations within the framework of Stochastic Geometry.

Introduction

- ▶ In recent years, it became popular to model the irregularity of the network base station (BS) locations within the framework of Stochastic Geometry.
- ▶ **Cooperation in Cellular Networks:** Very important subject of study. Why? Two or more BSs exchange user information to offer stronger signal with reduced interference.
 - ▶ Coverage improvement.
 - ▶ Better service to **cell-edge** users (inter-cell interference).

Introduction

Benefit depends on:

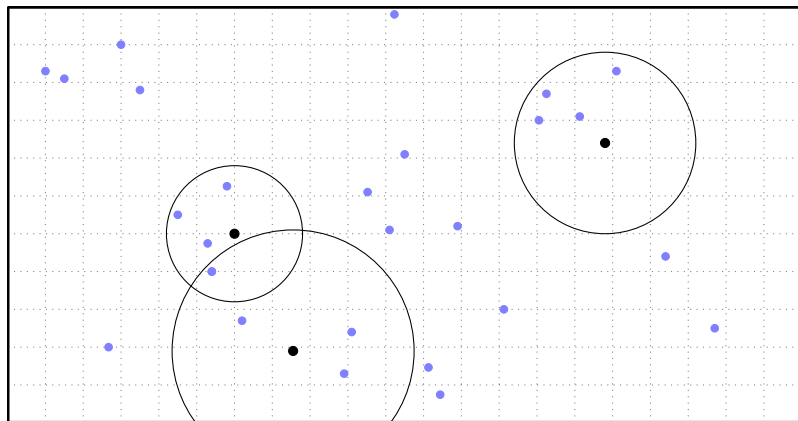
- ☞ Cooperation type.
- ☞ Information exchange.
- ☞ **Number and position of cooperating nodes!**

Literature

Stochastic Geometry (PPP) models:

- ▶ Andrews, Baccelli & Ganti (TCOM 2011): Tractable Approach
- ▶ Dhillon, Ganti, Baccelli & Andrews (JSAC 2012): K-tier HetNets
- ▶ Akoum & Heath (TSP 2013): Random layered clustering and cooperation.
- ▶ Baccelli & Giovanidis (Asilomar 2013 - TWC 2015): Cooperation between base station pairs.
- ▶ Nigam, Minero & Haenggi (TCOM 2014): Cooperation in groups of more than two.
- ▶ Tanbourgi, Singh, Andrews & Jondral (TWC 2014): Circle of fixed radius around user.
- ▶ Keeler, Błaszczyszyn & Karray (ISIT 2013, arXiv 2014): k-coverage, factorial moment measures.

Dynamic clusters: User-Driven



Every user (black dots) is served by its four closest BSs (blue dots).

Static Vs Dynamic

Problems with Dynamic Clusters:

- ▶ Each user can ask **any** set of BSs to cooperate → overburdens the backhaul/control channel with intensive communication.
- ▶ Time- (or more general Resource-) sharing between clusters.

Static Vs Dynamic

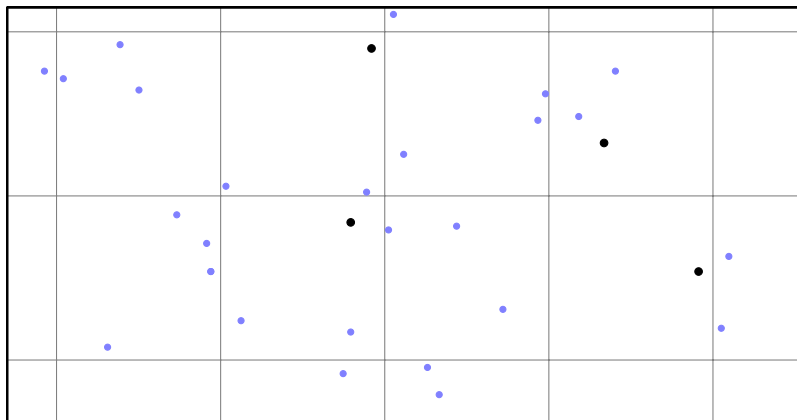
Problems with Dynamic Clusters:

- ▶ Each user can ask **any** set of BSs to cooperate → overburdens the backhaul/control channel with intensive communication.
- ▶ Time- (or more general Resource-) sharing between clusters.

Could **static** cooperation groups be more realistic? **YES!**

- ▶ Reasonable information exchange between BSs.
- ▶ No resource-sharing.
- ▶ One controller per cluster: Better to program and coordinate.

Static Clusters: Randomized & Hierarchical



Clusters are formed by those BSs (blue dots) inside each square.

Static Clusters: Proximity

👉 Define a-priori static groups by means of **proximity**.

👉 Why?

- ▶ Strong signal.
- ▶ Weak interference (avoid first-order).
- ▶ Fast coordination.
- ▶ BSs share a planar area of common interest.

Here: Singles and Pairs

The maximum group size constraint is $K = 2$:

- ▶ **single** BSs operating individually,
- ▶ **pairs** of BSs providing service cooperatively.

Here: Singles and Pairs

The maximum group size constraint is $K = 2$:

- ▶ **single** BSs operating individually,
- ▶ **pairs** of BSs providing service cooperatively.

Why?

- ▶ To be avoided: big number of BSs serving only one user.
- ▶ Easier to derive first results.

Mutually Nearest Neighbour Pair

Single

The point processes $\Phi^{(1)}$ and $\Phi^{(2)}$

Definition

Fixed a realisation ϕ ,

$$\phi^{(1)} = \{x \in \phi \ \& \ x \text{ is single } \},$$

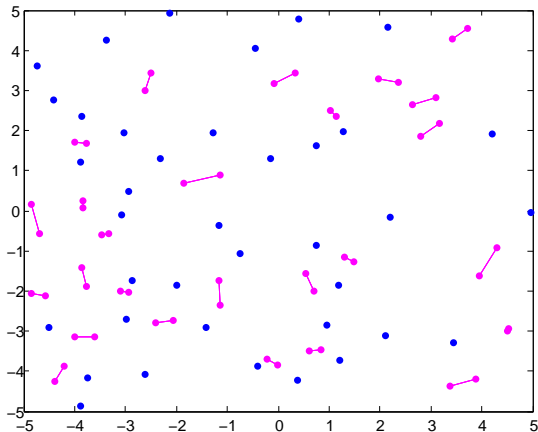
$$\phi^{(2)} = \{x \in \phi \ \& \ x \text{ cooperates with another element of } \phi \}.$$

Given a stationary point process $\Phi = \{\phi\}$ we define the point process $\Phi^{(1)}$ and $\Phi^{(2)}$ by means of the **dependent thinning** defined above:

$$\Phi^{(1)} = \{\phi^{(1)}\}$$

$$\Phi^{(2)} = \{\phi^{(2)}\}.$$

The point processes $\Phi^{(1)}$ and $\Phi^{(2)}$



Some results.

If Φ is an homogeneous PPP,

- ▶ 62% of points are pairs, and 38% are single points (it is impossible for $\Phi^{(1)}$ and $\Phi^{(2)}$ to be PPP).
- ▶ A representations for the Palm measures of $\Phi^{(1)}$ and $\Phi^{(2)}$.
- ▶ A closed formula for the NN function of $\Phi^{(2)}$.

Some numerical results:

- ▶ **Repulsion** between the process of singles, and **attraction** between the process of pairs.
- ▶ Tight bounds for the **ES** functions of $\Phi^{(1)}$ and $\Phi^{(2)}$.
- ▶ Average Voronoi surface proportions.

Interference field

Interference field generated by $\Phi^{(1)}$ and $\Phi^{(2)}$:

$$\begin{aligned} \mathcal{I}^{(1)} &= \sum_{x \in \Phi^{(1)}} f(x), \\ \mathcal{I}^{(2)} &= \sum_{x \in \Phi^{(2)}} \sum_{y \in \Phi^{(2)}} \frac{1}{2} g(x, y) \mathbf{1}_{\{x \leftrightarrow y\}}. \end{aligned}$$

Expected Interference

Theorem

The expected value of the interference fields has the analytical representation

$$\mathbb{E}[\mathcal{I}^{(1)}] = \int_{\mathbb{R}^2} \mathbb{E}[f(x)](1 - p^*)\lambda dx,$$
$$\mathbb{E}[\mathcal{I}^{(2)}] = \int_{\mathbb{R}^2} \int_{\mathbb{R}^2} \frac{1}{2} \mathbb{E}[g(x, y)] e^{-\lambda|x-y|^2(2\pi-\gamma)} \lambda dy \lambda dx.$$

Laplace Transform for Singles

Theorem

For a random field $f : \mathbb{R}^2 \rightarrow \mathbb{R}^+$, and $A \subset \mathbb{R}^2$. Let $F^{(n)}(x_1, \dots, x_n) = (f(x_1), \dots, f(x_n))$, the LT of the interference $\mathcal{I}^{(1)}$ generated by f and $\Phi_A^{(1)}$ is equal to

$$\mathbb{E} \left[e^{-s\mathcal{I}^{(1)}} \right] = e^{-\lambda \mathcal{S}(A)} \left(1 + \lambda \int_A \mathbb{E} \left[e^{-sf(x)} \right] dx + \frac{\lambda^2}{2} + \sum_{n=3}^{\infty} \frac{\lambda^n}{n!} \int_A \dots \int_A \mathbb{E} \left[e^{-s(F^{(n)} \cdot H^{(n)})(x_1, \dots, x_n)} \right] dx_1 \dots dx_n \right).$$

Laplace Transform for Doubles

Theorem

For a random field $g : \mathbb{R}^2 \times \mathbb{R}^2 \rightarrow \mathbb{R}^+$, and $A \subset \mathbb{R}^2$. Let $G_i^{(n)}(x_1, \dots, x_n) = (g(x_i, x_1), \dots, g(x_i, x_n))$, and $G^{(n)} = (G_1^{(n)}, \dots, G_n^{(n)})$, the LT of the interference $\mathcal{I}^{(2)}$ generated by g and $\Phi_A^{(2)}$ is equal to

$$\begin{aligned} \mathbb{E} \left[e^{-s\mathcal{I}^{(2)}} \right] &= e^{-\lambda \mathcal{S}(A)} \left(1 + \lambda \mathcal{S}(A) \right. \\ &\quad + \frac{\lambda^2}{2} \int_A \int_A \mathbb{E} \left[e^{-\frac{s}{2}(g(x,y)+g(y,x))} \right] \lambda dy \lambda dx \\ &\quad \left. + \sum_{n=3}^{\infty} \frac{\lambda^n}{n!} \int_A \dots \int_A \mathbb{E} \left[e^{-\frac{s}{2}(G^{(n)} \cdot J^{(n)})(x_1, \dots, x_n)} \right] dx_1 \dots dx_n \right). \end{aligned}$$

Extensions.

- ▶ Higher moments of $\Phi^{(1)}$ and $\Phi^{(2)}$.
- ▶ A rigorous proof for the numerically results.
- ▶ Use the above results to analyze the performance of cooperating cellular networks.
- ▶ Cooperation functions.

Extensions

$(v_x)_{x \in \mathbb{R}^2}$ random propagation effects from the BS x to the typical user.

- ▶ No cooperation function, $f(x) = \frac{v_x}{\|x\|^\beta}$.
- ▶ Pair cooperation function,

$$g(x, y) = \begin{cases} \frac{v_x}{\|x\|^\beta} + \frac{v_y}{\|y\|^\beta}, & \text{[NC]} \\ \max \left\{ \frac{v_x}{\|x\|^\beta}, \frac{v_y}{\|y\|^\beta} \right\}, & \text{[OF1]} \\ \mathbb{1}_{on_x} \frac{v_x}{\|x\|^\beta} + (1 - \mathbb{1}_{on_x}) \frac{v_y}{\|y\|^\beta}, & \text{[OF2]} \\ \left| \sqrt{\frac{v_x}{\|x\|^\beta}} e^{i\theta(x)} + \sqrt{\frac{v_y}{\|y\|^\beta}} e^{i\theta(y)} \right|^2 & \text{[PH]}. \end{cases}$$

What kind of laws are we interested in?

Extensions

- ▶ Because of the difficulties in analysing the previous schemes, we have further proposed a model that imitates the cluster structure of the BSs:
 - ▶ Based on using a superposition of 2 Poisson processes.
 - ▶ One is for the singles.
 - ▶ The second for the pairs.
 - ▶ The percentage of single and pairs.
 - ▶ The distance between two cooperating BSs.

We have obtained closed, analytic formulas.

Thank you for your attention.

 luis.alvarez-corrales@telecom-paristech.fr